



# How To Identify Hidden Markov Models

## An Algebraic Statistical Answer

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# Guideline

## Introduction

- Hidden Markov Chains
- Identifiability

## Algebraic Statistics

- The Hidden Markov Model
- The Matrix Markov Model
- Differences

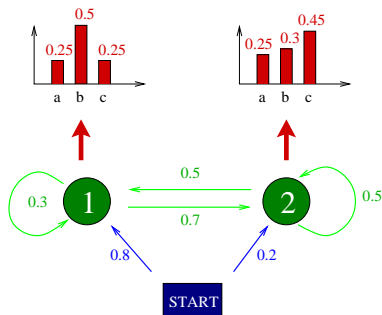
## The Main Theorem

- The Hankel Matrix
- Invariants

## Outlook



# Hidden Markov Chains



- Initial probabilities  $\pi = (0.8, 0.2)^T$
- Transition probabilities

$$M = (m_{ij} := P(j \rightarrow i))_{i,j=1,2} \\ = \begin{pmatrix} 0.3 & 0.5 \\ 0.7 & 0.5 \end{pmatrix}$$

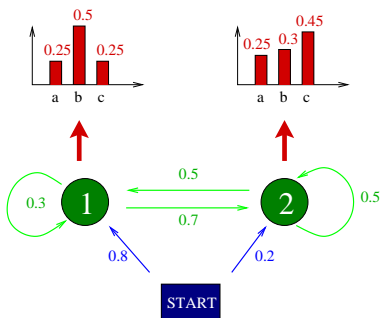
- Emission probabilities, e.g.  $E_{1b} = 0.5$ ,  $E_{2c} = 0.45$ .

Random source  $(X_t)$  with values in  $\Sigma = \{a, b, c\}$ :

$$\text{e.g.: } P_X(X_1 = a, X_2 = b) = \pi_1 e_{1a}(a_{11} e_{1b} + a_{12} e_{2b}) + \pi_2 e_{2a}(a_{21} e_{1b} + a_{22} e_{2b})$$



## Alternative Formulation



- Initial probabilities  $\pi = (0.8, 0.2)^T$
- Transition probabilities

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- Emission probabilities, e.g.  $E_{1b} = 0.5$ ,  $E_{2c} = 0.45$ .

Alternative formulation: Let

$$T_v := M \begin{pmatrix} E_{1v} & 0 \\ 0 & E_{2v} \end{pmatrix}, v = a, b, c$$

then

$$P_X(X_1 = v_1, \dots, X_n = v_n) = \langle \begin{pmatrix} 1 \\ 1 \end{pmatrix} | T_{v_n} \cdot \dots \cdot T_{v_1} | \begin{pmatrix} \pi_1 \\ \pi_2 \end{pmatrix} \rangle.$$



# Identifiability of HMMs

[Blackwell and Koopmans, 1957]

- **Identifiability Problem:** When do two HMMs give rise to equivalent processes?

Solutions:

- Ito, Amari, Kobayashi (1992): exponential runtime
- AS (2008): linear runtime
- When is a process due to an HMM? Solution:
  - Heller (1965): characterization of no practical use
- **Remark:** Hidden Markov process on  $d$  hidden states is uniquely determined by its distribution on strings of length  $2d - 1$ .



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- **Remark:** Hidden Markov process on  $d$  hidden states is uniquely determined by its distribution on strings of length  $2d - 1$ .
- Question of practical relevance: when is a distribution on strings of length  $n$  induced by an HMM with at most  $\lceil \frac{n}{2} \rceil$  hidden states?
- If so, how many are there?
- Algebraic statistics can give answers!



# Algebraic Statistical Models

## The Hidden Markov Model

**Definition:** **Hidden Markov Model** for  $d$  hidden states and strings of length  $n$  over the  $\Sigma$ :

$$\mathbf{f}_{n,d} : \mathbb{C}^{d(d-1)+d|\Sigma|} \longrightarrow \mathbb{C}^{|\Sigma|^n}$$

$$((T_a = MO_a)_{a \in \Sigma}, \pi) \mapsto (\langle \mathbf{1} | T_{v_n} \cdot \dots \cdot T_{v_1} | \pi \rangle)_{v=v_1 \dots v_n \in \Sigma^n}$$

where  $\pi, \mathbf{1} = (1, \dots, 1)^T \in \mathbb{C}^d$ ,  $M, O_a \in \mathbb{C}^{d^2}$  such that

$$\mathbf{1}^T M = \mathbf{1}^T$$

$$(O_a)_{ij} = \begin{cases} E_{ia} & i = j \\ 0 & i \neq j \end{cases}, \quad \sum_a O_a = \text{Id}.$$

**Wanted:** Equations describing

$$\overline{\text{im}(\mathbf{f}_{n,d})}$$



# Algebraic Statistical Models

## The Matrix Markov Model

**Definition:** Matrix Markov Model of rank  $d$  for strings of length  $n$

$$\mathbf{g}_{n,d} : \begin{array}{ccc} \mathbb{C}^{|\Sigma|d^2} & \longrightarrow & \mathbb{C}^{|\Sigma|^n} \\ ((T_a)_{a \in \Sigma}, \pi) & \longmapsto & (\langle \mathbf{1} | T_{v_n} \dots T_{v_1} | \pi \rangle)_{v=v_1 \dots v_n \in \Sigma^n} \end{array}$$

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where  $\mathbf{1}^T (\sum_a T_a) = \mathbf{1}$ .

**Reminder:** **Hidden Markov Model** on  **$d$  hidden states** for **strings of length  $n$**

$$\mathbf{f}_{n,d} : \begin{array}{ccc} \mathbb{C}^{d^2 + (|\Sigma|-1)d} & \longrightarrow & \mathbb{C}^{|\Sigma|^n} \\ ((T_a = MO_a)_{a \in \Sigma}, \pi) & \mapsto & (\langle \mathbf{1} | T_{v_n} \dots T_{v_1} | \pi \rangle)_{v=v_1 \dots v_n \in \Sigma^n} \end{array}$$

Hence

$$\text{im}(\mathbf{f}_{n,d}) \subset \text{im}(\mathbf{g}_{n,d}).$$



# Matrix Markov Model

## Dimension

**Lemma:** Let  $S : \mathbb{C}^d \rightarrow \mathbb{C}^d$  be an **isomorphism** such that  $\mathbf{1}^T S = \mathbf{1}^T$  and

$$T'_a := S^{-1} T_a S$$

$$\mathbf{x}' := S^{-1} \mathbf{x}.$$

Then

$$\mathbf{g}_{n,d}((T_a)_{a \in \Sigma}, \mathbf{x}) = \mathbf{g}_{n,d}((T'_a)_{a \in \Sigma}, \mathbf{x}')$$

**Theorem:** Let  $n \geq 2d - 1$ .

$$\dim \overline{\text{im } \mathbf{g}_{n,d}} = (|\Sigma| - 1)d^2 + d.$$



# Differences

## The Strassen Condition

**Lemma:** Equivalent statements:

(i)

$$\mathbf{g}_{n,d}((T_a)_{a \in \Sigma}, \mathbf{x}) \in \text{im } \mathbf{f}_{n,d}$$

(ii) [Strassen Condition]

$$\forall a, b, c \in \Sigma : (T_a)^{-1} T_c (T_b)^{-1} = (T_b)^{-1} T_c (T_a)^{-1}.$$

**Remark:** The Strassen Condition reflects that

$$T_c(T_a)^{-1}, T_c(T_b)^{-1}$$

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**Corollary:**

$$\Sigma = \{a, b\} : \overline{\text{im } \mathbf{f}_{n,d}} = \overline{\text{im } \mathbf{g}_{n,d}}$$

since (ii) trivially satisfied for binary alphabet  $\Sigma = \{a, b\}$ .



# Main Theorem

## The Hankel Matrix

- Let  $v \in \Sigma^n$ . Write

$$p(v = v_1 \dots v_n) \\ := \mathbb{P}(X_1 = v_1, \dots, X_n = v_n).$$

for stochastic process  $(X_t)$ .

- Write

$$uv = u_1 \dots u_m v_1 \dots v_n \in \Sigma^{m+n}$$

for **concatenation** of

$$u = u_1 \dots u_m \in \Sigma^m \text{ and}$$

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- Consider the (infinite-dimensional) **Hankel matrix**

$$\mathcal{P}_p := [p(uv)]_{u,v \in \Sigma^*} \in \mathbb{R}^{\Sigma^* \times \Sigma^*}.$$

- Example** ( $\Sigma = \{0, 1\}$ ):

$$\mathcal{P}_p = \begin{pmatrix} p(\square) & p(0) & p(1) & \dots \\ p(0) & p(00) & p(01) & \dots \\ p(1) & p(01) & p(11) & \dots \\ p(00) & p(000) & p(010) & \dots \\ p(01) & p(001) & p(011) & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$



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**Remark:** For  $p$  **hidden Markov process** on  $d$  hidden states:

$$\text{rk } \mathcal{P}_p \leq d.$$



# Main Theorem

## Set Theoretic Lemma

**Lemma:** Let  $n \geq 2d - 1$ . Then the following statements are equivalent:

(i)

$$(\rho(v))_{v \in \Sigma^n} \in (\text{im } \mathbf{g}_{n,d} \setminus \text{im } \mathbf{g}_{n,d-1}) \quad (1)$$

(ii)

$$\text{rk } \mathcal{P}_{\rho, d-1, d-1} = \text{rk } \mathcal{P}_{\rho, \lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil} = \mathcal{P}_{\rho, \lceil \frac{n}{2} \rceil, \lfloor \frac{n}{2} \rfloor} = d \quad (2)$$





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(ii)

$$\text{rk } \mathcal{P}_{p,d-1,d-1} = \text{rk } \mathcal{P}_{p, \lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil} = \mathcal{P}_{p, \lceil \frac{n}{2} \rceil, \lfloor \frac{n}{2} \rfloor} = d \quad (2)$$

**Example:**  $n = 4, d = 2$

$$\mathcal{P}_{p,4,2} = \begin{pmatrix} p(\square) & p(0) & p(1) & p(00) & p(01) & p(10) & p(11) \\ p(0) & p(00) & p(10) & p(000) & p(010) & p(100) & p(110) \\ p(1) & p(01) & p(11) & p(001) & p(011) & p(101) & p(111) \\ p(00) & p(000) & p(100) & p(0000) & p(0100) & p(1000) & p(1100) \\ p(01) & p(001) & p(101) & p(0001) & p(0101) & p(1001) & p(1101) \\ p(10) & p(010) & p(110) & p(0010) & p(0110) & p(1010) & p(1110) \\ p(11) & p(011) & p(111) & p(0011) & p(0111) & p(1011) & p(1111) \end{pmatrix}$$



# Main Theorem

## Invariants

Let

$$I := \langle \det(p(u_i v_j))_{1 \leq i, j \leq d+1} \mid 0 \leq |u_i|, |v_j| \leq \lceil \frac{n}{2} \rceil, |u_i v_j| \leq n \rangle$$

$$J := \langle \det(p(u_i v_j))_{1 \leq i, j \leq d} \mid 0 \leq |u_i|, |v_j| \leq d-1 \rangle.$$

**Note:** To obtain equations for strings of length  $n$ , replace

$$p(uv) = \sum_{w, |w|=n-|uv|} p(uvw) \quad \text{for } |uv| < n.$$

**Theorem:**

$$\overline{\text{im } \mathbf{g}_{n,d}} = \text{rad } I : \text{rad } J.$$



# Outlook

- Larger alphabets
- Pathological cases
- Transporter theorem



***Thanks for the attention!***